

Blockchain Interoperability

Goals of blockchain

- ▶ Borderless
- ▶ Censorship resistant (neutrality)
- ▶ Open

Applications

- ▶ Healthcare ¹
- ▶ Finance
- ▶ Supply-chain ²

¹Holbl, Marko, Kompara, Marko & Nemeč Zlatolas, Lili (2018) 'A systematic review of the use of blockchain in healthcare', *Symmetry* 10, 470

²Notheisen, Benedikt & Shanmugam, Arun Prasad (2017) 'Trading real-world assets on blockchain', *Business & Information Systems Engineering*

Prevalent current situation

- ▶ Every blockchain can't communicate with other blockchains
- ▶ No extension

Consequences:

- ▶ Vendor lock-in and centralization
- ▶ Complex blockchain design and less security
- ▶ Coin Maximalism
- ▶ No blockchain adoption

Interoperability advantages

- ▶ Help solving scalability vs decentralization trade-off
- ▶ Add blockchain features
- ▶ Different assets besides currencies

Three Layers Structure

The three layers of abstraction:

- ▶ Layer 0: message exchange
- ▶ Layer 1: blockchains structure
- ▶ Layer 2: side-channels

Layer 0 - Basic Idea

- ▶ Clients of a blockchain can modify only their blockchain
- ▶ Clients of a blockchain can view changes of other blockchains

Layer 0 - Methods

- ▶ Proof of Lock ³
- ▶ Notary schemes ⁴
- ▶ Atomic Swaps ⁵

³Back, Adam, Corallo, Matt, Dashjr, Luke, Friedenbach, Mark, Maxwell, Gregory, Miller, Andrew, Poelstra, Andrew & Wuille, Pieter (2014) 'Enabling Blockchain Innovations with Pegged Sidechains'

⁴Dilley, Johnny, Poelstra, Andrew, Wilkins, Jonathan, Piekarska, Marta & Friedenbach, Mark (2016) 'Strong federations: An interoperable blockchain solution to centralized third-party risks', *arXiv preprint arXiv:1612.05491*

⁵Miraz, Mahdi & Donald, David C (2019) 'Atomic Cross-chain Swaps: Development, Trajectory and Potential of Non-monetary Digital Token Swap Facilities', *Annals of Emerging Technologies in Computing (AETiC) Vol 3*

Proof of Lock

1. Client A locks funds on blockchain B_1
2. Client A has/produces proof of lock and presents it to clients on blockchain B_2
3. A client B on blockchain B_2 verifies the proof and releases/creates equivalent funds on blockchain B_2

Proof of Lock - Comments

Con:

- ▶ Implements light-clients
- ▶ Risk of blocking transaction
- ▶ Slow
- ▶ Need to peg blockchain

Pro:

- ▶ Easy to implement
- ▶ Decentralized

Notary schemes

- ▶ Third party does the exchange
- ▶ Examples
 - ▶ Online exchanges
 - ▶ Smart contract
 - ▶ Federation with majority

Notary schemes - Comments

Con:

- ▶ Centralized
- ▶ Single point of failure

Pro:

- ▶ Easy to implement
- ▶ Fast

Atomic Swaps

- ▶ Either successfully concluded or revert
- ▶ Currently only with hashtimelock commitment schemes
 - ▶ Startup phase
 - ▶ Commitment phase
 - ▶ Claim phase

Hashtimelock - Startup phase/1

1. Alice creates a number x , hash it $H(x)$ and gives the hash to Bob)
2. Alice creates a transaction T_1 which would transfer on B_1 n_1 tokens from Alice's address to Bob's one
 - ▶ Can be spent only providing x via a message signed by Bob or with a general transaction signed by both Alice and Bob

Hashtimelock - Startup phase/2

3. Alice creates a transaction T_2 which would transfer, after t hours, on B_1 the tokens from T_1 to her address; she sends this transaction to Bob
4. Bob signs transaction T_2 and returns it to Alice
 - ▶ This is Alice's warranty that even if Bob turns out to be malicious, she will have her tokens back

End of the startup phase

Hashtimelock - Commitment phase/1

1. Alice publishes T_1 on B_1
2. Bob creates a transaction T_3 which would transfer on B_2 n_2 tokens from Bob's address to Alice's one
 - ▶ Can be spent only providing x via a message signed by Alice or with a general transaction signed by both Alice and Bob

Hashtimelock - Commitment phase/2

3. Bob creates a transaction T_4 which would transfer, after $\frac{t}{2}$ hours, on B_2 the tokens from T_3 to his address; he sends this transaction to Alice
4. Alice signs transaction T_4 and returns it to Bob
 - ▶ This is Bob's warranty that even if Alice turns out to be malicious, he will have his tokens back
5. Bob publishes T_3 on B_2

End of commitment phase

Hashtimelock - Claim phase

1. Alice spends T_3 before $\frac{t}{2}$ hours and gives x
2. Bob spends T_1 before t hours using x

Hashtimelock - Comments

Con:

- ▶ Difficult to implement
- ▶ Slow
- ▶ Prone to DDoS
- ▶ Weak in case of reorg
- ▶ Same hash function on both blockchain

Pro:

- ▶ Censorship resistant
- ▶ Highest decentralization

Layer 1 - Configurations

- ▶ Tree
- ▶ Star
- ▶ Parallel chains ⁶

⁶Lerner, Sergio Demian (2016) *Drivechains, Sidechains and Hybrid 2-way Peg Designs_R9*

Tree

- ▶ Root chain
- ▶ Every blockchain reports its state to the parent chain for inter-chain transaction
- ▶ Examples:
 - ▶ Treechain (deprecated)
 - ▶ Plasma ⁷

⁷Poon, Joseph & Buterin, Vitalik () 'Plasma: Scalable Autonomous Smart Contracts'

Star

- ▶ One central blockchain
- ▶ All the blockchain communication must pass by the central hub
- ▶ Example:
 - ▶ Cosmos
 - ▶ Polkadot ⁸

⁸Wood, Dr Gavin 'Polkadot: Vision for a Heterogeneous Multi-chain framework'

Parallel chains

- ▶ Separated blockchains capable of exchanging messages
- ▶ Examples:
 - ▶ RSK
 - ▶ Blockstack

Layer 2 - Side-channels

- ▶ Not all the transaction must be recorded (e.g. smaller recurrent payments)
- ▶ Active development because of scalability improvements
- ▶ Examples:
 - ▶ Lightning Network (Bitcoin)
 - ▶ Raiden Network (Ethereum)

Contacts

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